On the intersection of Information Centric Networking and Delay Tolerant Networking (Lessons learned from the GreenICN project)

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# Outline

- (Very) Short Introduction to Information Centric Networking
- GreenICN Project Background
  - Disaster Scenario
  - Rationale
- Overview on selected Solutions
  - ...which we developed and evaluated
- Deep Dive
  - Decentralised ICN Interest Popularity Estimation
- Lessons Learned & Open Questions
  - Key Takeaways
  - Remaining Issues & Challenges
  - Discussion

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### (Very) Short Introduction to Information Centric Networking

# **Host-centric networking**



# **Information-Centric Networking**



# **ICN communication model**

- Clients (C) send requests (Interest Packets) asking for named data
- Routers (R) in the network route requests towards publishers (P)
- Any node with a cached copy can provide the corresponding information object (Data Packet)

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- Pending Interest Table (PIT): "stores all the Interests that a router has forwarded but not satisfied yet" [https://en.wikipedia.org/wiki/Named\_data\_networking]
- Remark:
  - On the surface, this is exactly the service of HTTP, but the request is there always addressed to a particular host



# **ICN-based Information Retrieval 101**



# **ICN Core Properties**

- Accessing named data as a first-order network principle
  - Transmission of self-contained units
- Name-content-binding validation and other security services based on object/naming security
  - Not based on connection security

### Ability to leverage ubiquitous in-network memory

- Rate adaptation
- Repair (efficient re-transmissions)
- Sharing (Re-use)

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### **GreenICN Project Background**

# <sup>97.12.19</sup> Background: GreenICN Project

- GreenICN: Architecture and Applications of Green Information Centric Networking
- Duration: 3,16 years (1 Apr 2013 31 May 2016)
- Website: <u>http://www.greenicn.org</u>

EU Coordinator: Prof. Xiaoming Fu University of Göttingen Germany JP Coordinator: Mr. Shigehiro Ano KDDI R&D Labs Japan



### 17.12.19

# **Project Consortium**

#### European Partners



GEORG-AUGUST-UNIVERSITÄT GÖTTINGEN

EU Coordinator Georg-August-Universität Göttingen (UGO, Germany) Contact: Xiaoming Fu <fu@cs.uni-goettingen.de>



CEDEO (CED, Italy)

orange"

Telekomunikacja Polska (Orange Labs, Poland)

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#### Japanese Partners



JP Coordinator KDDI R&D Laboratories Inc. (KDD, Saitama) Contact: Shigehiro Ano <ano@kddilabs.jp>

NEC Corporation (NEJ, Tokyo)



Panasonic Advanced Technology Development Co., Ltd







### Scenario and Use Cases

### **Disaster Scenario**

- The aftermath of a disaster (hurricane, earthquake, tsunami, or a human-generated network breakdown)
- E.g. the enormous earthquake which hit Northeastern Japan on March 11, 2011 (causing extensive damages incl. blackouts, fires, tsunamis and a nuclear crisis)



- Energy and communication resources are at a premium
- Critical to efficiently distribute disaster notification and critical rescue information

### Key Use Cases (High-Level)

- Authorities would like to inform citizens of possible shelters, food, or of impending danger
- Relatives would like to communicate with each other and be informed about their wellbeing
- Affected citizens would like to make enquiries of food distribution centres, shelters or report trapped, missing people to the authorities

Key Research Challenges

- Communication in Fragmented Networks (using disconnected but functional parts of the infrastructure)
- Security (access control, message authentication)
- Traffic Prioritization / Handling Congestion (overall capacity is reduced)
- Delay/Disruption Tolerant Approach



# How ICN can be Beneficial<sup>[1,2]</sup>

### Routing-by-Name

• In fragmented networks, references to location-based, fixed addresses may not work as a consequence of disruptions (e.g. reachability of DNS servers)

### Content-based Access Control

• ICN security model can regulate access to data objects (e.g. only to a specific user or class of users) by means of content-based security

### Authentication of Named Data Objects

• With 'self-certifying data' approaches, the origin of data retrieved from the network can be authenticated without relying on a trusted third party or PKI

### Caching

 Caching can help to avoid congestion in the network (e.g. congestion in backhaul links can be avoided by delivering content from caches at access nodes)

### Sessionless Communication

• ICN does not require full end-to-end connectivity (facilitating a seemless aggregation between normal operations and a disaster)

 [1] J. Seedorf et al.: "Using ICN in disaster scenarios", draft-irtf-icnrg-disaster-09, IRTF ICNRG, Dec. 2019
 [2] J. Seedorf et al.: "The Benefit of Information Centric Networking for Enabling Communications in Disaster Scenarios", Globecom 2015 Workshop on Information Centric Networking Solutions for Real World Applications (ICNSRA), San Diego, USA, December, 2015

# Research Gap

- Quite some work in the DTN community, however most DTN work lacks key features which are needed in the disaster scenarios we consider, such as:
  - publish/subscribe (pub/sub) capabilities, caching, multicast delivery, message prioritisation based on content types, ...
- Could enhance existing DTN approaches with these features we argue that ICN makes a better starting point for building a communication architecture that works well before & after a disaster

# → Vision / Rationale: Start with existing ICN approaches and extend them with the necessary features needed in disaster scenarios

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### **Overview on selected Solutions**

- ICN 'Data Mules' [2] [3]
  - Logical interface, multipath support



ICN Data Mules in a Disaster Scenario

 [2] T. Yagyu and S. Maeda, "Demo Overview: Reliable Contents Retrieval in Fragmented ICNs for Disaster Scenario," ACM ICN Conf., Sep. 2014.
 [3] K. Sugiyama et al., "Multipath Support for Name-based Information Dissemination in Fragmented Networks," ACM ICN Conf., Sep./Oct. 2015. 17.12.19

- ICN 'Data Mules'
- Priority dependent Name-based Replication (NREP) [4]
  - Routing/forwarding decisions based on name/attributes
  - E.g. attaching priority & time/space restrictions to interests



More Replications till Expiry for High Priority Messages

[4] I. Psaras et al., "Name-based replication priorities in disaster cases," in 2<sup>nd</sup> Workshop on Name Oriented Mobility NOM), 2014.

- ICN 'Data Mules
- Priority dependent Name-based Replication (NREP)
- Data-centric Confidentiality/Access Control/Authentication [5]
  - Multi-authority 'Ciphertext-Policy Attribute Based Encryption' ICN security architecture
  - Example Policy: allow access only to recipients who fulfill:

 $\Pi = (job:official \land rank:executive) \lor (job:emergency \land rank:any)$ 



[5] T. Asami et al., "D2.3.1 - initial solution for access control and management in fragmented networks," GreenICN Project, GreenICN Project Deliverable, 2013

- ICN 'Data Mules'
- Priority dependent Name-based Replication (NREP)
- Data-centric Confidentiality/Access Control/Authentication [6][7]
  - Binding between self-certifying ICN names and Real-World Identities via a Web-of-Trust (WoT)



[6] J. Seedorf, D. Kutscher, and F. Schneider: "Decentralised binding of self-certifying names to real-world identities for assessment of third-party messages in fragmented mobile networks," 2nd Workshop on Name Oriented Mobility (NOM), 2014
 [7] J. Seedorf et al.: "Demo overview: Fully decentralised authentication scheme for icn in disaster scenarios (demonstration on mobile terminals)."

in 1st ACM Conference on Information-Centric Networking (ICN-2014), 2014.

#### 17.12.19 [\*] https://github.com/BenjaminSchiller/GTNA/tree/master/src/gtna

- ICN 'Data Mules'
- Priority dependent Name-based Replication (NREP)
- Data-centric Confidentiality, Access Control, Authentication
- Information Resilience [8]
  - NDN Extension "Satisfied Interest Table (SIT)"
  - Keeps track of data packet next hop, i.e. stores info reg. downstream delivery of content
  - Use of this additional routing table in case upstream routing of interests provides no data



[8] V. Sourlas et al. "Information Resilience through User-Assisted Caching in Disruptive Content-Centric Networks", IFIP Networking 2015, May 2015 17.12.19



# Information Resilience through SIT



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### Deep Dive: Decentralised ICN Interest Popularity Estimation

# **Background & Scope**

#### **Background: Disaster Scenario**

- GreenICN Project: Using ICN for Disaster Scenarios
- Nodes may be scattered across different fragmented networks
- DTN-like communication using ICN data mules

#### **Goal: Decentralised Popularity Estimation of Interest Messages**

- Scenario: ICN data mules with limited storage capabilities and possibly limited delivery time
- Assumption: random, unpredictable movements of ICN data mules
- Date mules need to prioritize ICN messages
- Knowing interest popularity can be an important factor in prioritization



# **Proposed Solution: Rationale & Overview**

#### **General Idea**

- Exploit Nonces in ICN Protocols for counting interests received
- Naïve Approach: Append to each end-user request a unique nonce
  - Data mules would need to keep all nonces received in order not to over-count when encountering again the same end user or data mule in the future
  - Accurate estimation, but clearly does not scale ...

### **Overview of our Proposed Solution**

- Idea: Aggregate [nonce:counter] tuples
  - Approximate content popularity
  - Much more scalable (with respect to memory requirements at nodes)
- General Scheme:
  - End-users assign random nonces to interests
  - Data mules maintain list of [nonce:counter] tuple per interest for scalability
  - When two data mules meet, they exchange their Interests (incl. [nonce:counter])
    - Interests & popularity estimation gets distributed in network

# **Proposed Solution: Rationale & Overview**

#### Algorithm

- When two data mules encounter each other and both have for a given name already a [nonce:counter] tuple, aggregation of nonces and counters is performed
- For each Interest, aggregate [nonce:counter] tuples as follows:

```
Compare([nonce1, count1] , [nonce2, count2])
IF nonce1 == nonce2
    new_count = MAX(count1, count2) (at both nodes)
IF nonce1 != nonce2
    New_nonce = nonce with largest counter([nonce1, count1], [nonce2, count2])
    New_count = count1 + count2
```

# **Proposed Solution: Loop Prevention**

#### Handling Repeated Data Mule Encounters over Time

- Memory list: Data mules keep list of recently encountered mules per Interest
- If a data mule is encountered again, counters are not added
- Instead, max-counter rules is applied at both sides:

```
nonce1 != nonce2 (AND recently met)
```

- New\_nonce = nonce with largest counter([nonce1, count1], [nonce2, count2])
- New\_count = MAX(count1, count2)

### **Limited Memory List Size**

- Memory list of encountered nodes has limited (configurable) size
  - Can trade off accuracy (in popularity-prediction) against memory requirements at nodes
- When memory list full always use max-counter rule:
   New\_count = MAX(count1, count2) (=treat as if recently met)
- Alternative: Sliding approach = FIFO list of encountered nodes

# **Proposed Solution: Algorithm**



### **Proposed Solution: Nonce Assignment Example**



### **Proposed Solution: Message Flow Example**



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# **Evaluation**

### Simulations

- Objective: Evaluate accuracy of distributed popularity estimation
- Reference Algorithm: Naive solution
  - Append to each end-user request a unique nonce
  - At data mule encounters a complete list of nonces is exchanged for every given name in the Pending Interest Table (PIT) of each data mule
- Scenario / Experiments
  - 2 Fragmented Communities with 100 distinct data objects (size: 8kb)
  - Zipf distribution for Interests with 0.8 < alpha < 0.9 (appr. 500 interests issued by users)
  - Phase 1
    - 3 different data mules in each of the fragmented communities, receiving interests from users at random
  - Phase 2
    - All 6 Data mules meet each other in a random manner and exchange interests (no disruption during intermeeting times)



# **Results: Reference Algorithm**



# **Results: Our Approach**

**Results for Decentralised Interest Counter Aggregation for ICN (DICAI)** 



# **Results: Memory Requirements**

Memory utilization for nonces for the top 5 prefixes popularity-wise (among all data mules, 15 encounters, 1 Byte per nonce)



# **Summary & Contribution**

### Contribution

- Design of a scalable, fully decentralised scheme for estimating the popularity of ICN interest messages
  - scenario with random, unpredictable movements of ICN data mules
  - very useful to optimize content dissemination after a disaster
- Algorithm described in detail and with concrete examples
- Evaluation
  - Analytical model of the storage overhead introduced by our approach
  - Simulations showing that our proposed solution provides sufficient accuracy in predicting the actual content popularity

### **Future Work**

- More complex simulations
- Actual Implementation (e.g. based on CCNx, NDN)

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### Lessons Learned & Open Questions

# ICN is indeed a good match for Disaster Scenarios

### Our work confirmed that ICN is a good starting point for enabling disaster aftermath communication

- ICN brings (intrinsicly) many features that are very useful for enabling communications after a disaster happened
  - E.g. naturally supports sessionless communication
- We extended existing ICN approaches with the necessary features needed in disaster scenarios, showing that it is possible to enhance ICN with DTN features
  - Decentralized Forwarding
  - Message Prioritization
  - Security



EU-Japan FP7 GreenICN project successfully completed with a grade of "EXCELLENT"



#### www.greenicn.org

# ICN is indeed a good match for Disaster Scenarios

#### However, semantics of certain components may change ...

- Routing/forwarding essentially changes to a store-carry-forward of interests over a long period of time
  - For some use cases, rather 'forwarding to anybody' instead of 'routing'
- Examples
  - Nonces: Can also be used for interest popularity estimation, and may be aggregated
  - Interest Lifetime: May become more of an application-layer validity-period
  - Faces: Meaning may become somewhat blurry
    - Single physical face (e.g. Wifi) may be used to forward interests to various different nodes over time, independently of their ID (i.e. forward interests to anybody met)
  - Content forwarding: May even be meaningful to exchange popular data items at data mule encounter without a corresponding interest at that point in time

# Different Solutions for different Scenarios

- Fixed, predictable movement of ICN Data Mules
  - E.g. forward interests based on logical face
- Semi-random movement of ICN Data Mules (e.g. known direction of rescue teams)
  - E.g. forward interest based on prioritization and/or time/space attributes
- Completely random movement of ICN Data Mules (e.g. end users walking around)
  - E.g. forward interests to anybody based on interest popularity estimation

# Security

#### **Disaster Use Case good match for ICN Data-centric Security**

- IBE / ABE very suitable for key use cases
- Can assume single trusted key generator (e.g. government, mobile operator) for many use cases
- But may be difficult to apply these approaches on a world-wide scale
  - E.g. scalability / performance of ABE depends on number of attributes
  - Single trusted entity as PKG ...

#### Web-of-Trust fits well the decentralised nature of disaster scenario

- Enables fully decentralised authentication of content
- Our solution is based on self-certifying names (name contains the hash of the public key of the owner of the name)
- WoT provides only trustworthiness for a given name
  - Trust content based on relationships in the WoT graph
- Fully decentralised, but more-or-less 'probabilistic' security

# Security

#### Name Assignment Responsibility / Ownership of Names

- For hash-based names need to know the hash
  - If name contains hash of content, need to know the correct hash
  - If name contains hash of public key (of the owner of the name), need to know the corresponding Real-World Identity or public key
- Brings up the larger questions of how user obtains names
  - Outside the ICN layer?
  - Via trusted Search Engine?
  - Lengthy discussions ...
- For many of our solutions, just assumed pre-configured or well-known names
  - E.g. within the disaster app on your smartphone
  - E.g. bootstrap via well-known name where other producers can post names under which new content is available (like an 'alert bulletin board')
- May be very use case specific

# Other Open Issues

### Selected Open Issues\*

- Specifying for each mechanism suggested to what exact extent ICN deployment in the network and at user equipment is required
- How to best use DTN and ICN approaches for an optimal overall combination of techniques?
- How do data-centric encryption schemes scale and perform in large- scale, realistic evaluations?
- How to properly disseminate authenticated object names to nodes (for decentralised integrity verification and authentication) before a disaster, or how to retrieve new authenticated object names by nodes during a disaster?

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Questions & Discussion ...