Authentication and Access Control to Resources in a RELOAD Overlay

AG-iNET Seminar

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October 19, 2010



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■ Tightly coupled peer-to-peer conferences using SIP

- No central server for conference lookup, signaling and media distribution
- Focus functionality distributed transparently among participants
- Use RELOAD for registering conference URI





An Overview of RELOAD

- REsource LOcation And Discovery
- P2P-Overlay protocol
- Based on a distributed hash table (DHT), e.g. Chord
- Highly extensible by defining other usages, e.g. SIP- & XMPP-Usage
- In P2PSIP used to replace the proxy- and registrar-servers





The SIP-Usage for RELOAD

- Goal: Replacing the proxy and registar functionality in a P2P fashion
- Registration done by storing contact information using the SIP-Registration kind in the overlay
- Lookup of a SIP-URI done in normal DHT-Fashion: lookup hashed URI
- Proxy replaced by providing method to establish transport connection between any two peers (including NAT)





Similar to SIP-Usage

- Register a conference URI to make it publicly accessible
- Focus peers register themselves as entry points to the conference
- Write permission to the conference registration kind needs to be delegated to focus peers





The RELOAD Security Model

- Security model based on certificates
- Certificates used to provide authentication among peers
- Usually certificates obtained from central enrollment server
- RELOAD provides security on three levels:
 - Connection Level: (D)TLS for peer connections
 - Message Level: all messages signed
 - Object Level: all stored objects signed by creator
- Access control policies limit (write) access to resources based on certificates
- Policies indirectly prevent peers from storing arbitrary amounts of data





RELOAD

- Access control policies limit who may write at a specific location
- All stored data must be signed
 - → Storer presents certificate with fields matching policy rules
- Base RELOAD defines 4 policies:
 - USER-MATCH
 - NODE-MATCH
 - 3 USER-NODE-MATCH
 - NODE-MULTIPLE
- Every resource kind specifies policy to use e.g. SIP-Registration: USER-NODE-MATCH
- None of these fitting for DisCo-Registration



Introduction



Requirements for DisCo-Registration

- Shared resource (Needs to be written by multiple peers)
- Only (some) members of the conference allowed to write
- Anybody may retrieve stored data
- Group members need to verify stored data
- Must still fit into general RELOAD security model e.g. prevent resource exhaustion from storing lots of data (resource exhaustion)

Straightforward solution:

Use one conference certificate for the whole group





Ask Enrollment Server

Simplest approach:

- Get new certificate from enrollment server for each conference
- Certificate User-Name contains conference URI
- Use USER-MATCH policy
- Distribute private key of the conference certificate among focus peers

Problem:

- Enrollment Server only supposed to be contacted when joining the DHT:
 - Overlay should stay fully functional without server
- → Need to get a certificate without enrollment server





Self Signed Certificates Don't Work

- Malicious peer could take over a conference:
 - Generate certificate for existing conference URI (sybil attack)
 - Unable to tell if certificate is from conference initiator or not
 - Malicious peer registers itself as focus for the conference
- Peer could generate multiple certificates and store arbitrary amounts of data at multiple positions in the overlay





- Conference initiator creates conf-certificate, signed with his private key
- Define new access control policy for use in DisCo-Registration: USER-CHAIN-MATCH
- Storing peer verifies certificate chain
- Distribute conference private key to all focus peers





Chained Certificates Need Restrictions

Remaining Problem:

 Malicious peers could still generate certificates for existing conferences

Solution:

 Restrict who can create a conference and certificate for a certain name

- Conference name must be correlated to initiator's name
- Restrict allowed conference
 URIs using pattern matching

Example:

URI pattern: *-conf-\$USER@\$DOMAIN
User Name: alice@example.com

Allowed:

XYZ-conf-alice@example.com pretty-conf-alice@example.com NOT allowed:

alice-conference@example.com



laining i robiem. Revocation

- Chance of conference certificate being compromised enhanced, because multiple peers have the private key
- Still no solution for certificate revocation

Simple workaround:

- Short certificate lifetimes (e.g. a couple of minutes?)
- Renew certificate accordingly
- Certificate can only be used as long as creator is in the conference (+ lifetime)
 - ightarrow probably makes sense because conference name is coupled to creator

Alternative:

- Publish Certificate Revocation List (CRL) in overlay
- Needs another kind definition...



Introduction



Conclusion and Outlook

Introduction

- Controlled access to shared resources not a trivial problem in P2P networks
- RELOAD is flexible enough to support fitting access control policies
 - Are there better ideas than pattern matching?
- Certificate revocation can be done, but incurs more overhead Is it really needed for simple ad-hoc conferences?





Ich bedanke mich für die Aufmerksamkeit.

Fragen?

