

Internet Group Communication: IP Multicasting

- Introduction
- Why to Talk in Groups?
- Aspects of Group Communication
- IP-Multicasting
- Addressing
- The Internet Group Protocol

Motivation

Current Situation: Use of the Internet penetrates new areas:

- Multimedia Information Services
- Infotainment
- Gaming (MMORPGs)
- Synchronous Network Information Services
- Group Communication Tools

⇒ **A Transport Infrastructure is needed for
Group Communication Services**

Requirements

- Must fit into current IP Infrastructure
- Location Transparency
- Efficiency in Data Distribution
- Utmost Independence of Specific Hardware
- Independent Open Standards
- Interoperability

Example: Video Streaming

- High Requirements of Bandwidth and Server Performance
- Continuous Data Streams
- Real-Time Synchronisation
- Global Distribution (Web-TV, IPTV, Video-Conferencing)

Issues:

- **Internet Group Communication Fundamentals**
Group Communication, Multicasting, Addressing, IGMP/MLD
- **Layer 2 Multicasting**
Local Networks, Address-Mapping, Framing, Discovery, ATM
- **Multicast Routing**
Specialities, Algorithms, Protocols
- **New Developments**
IPng, SSM, Multicast Mobility

Why to Talk in Groups?

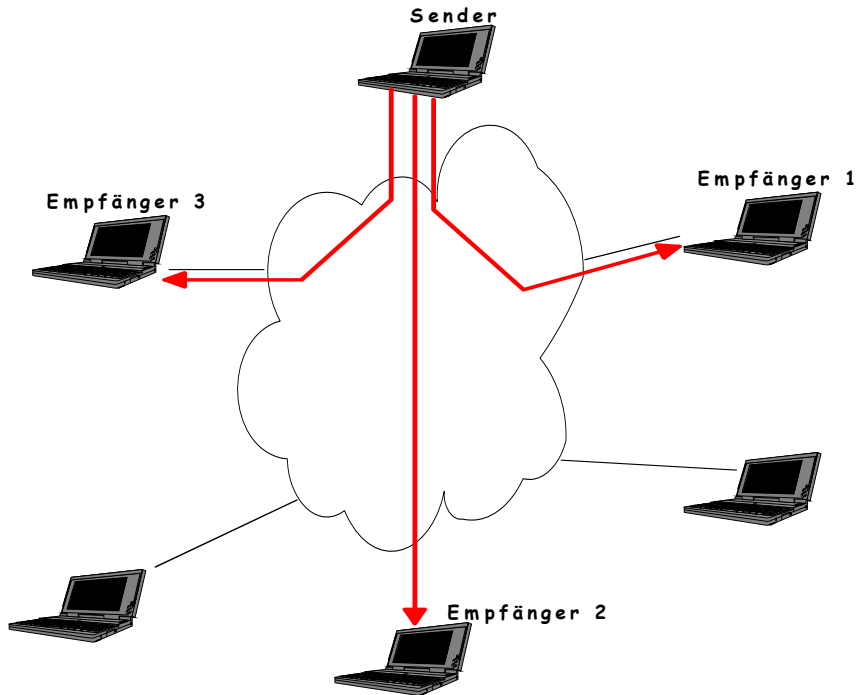
Internet based communication steadily gains importance, quantitatively as well as qualitatively. New communication forms arise, old services spread rapidly:

- Multimedia Distribution
- ‚Broadcasting‘ - Offers
- Telecommunication Services

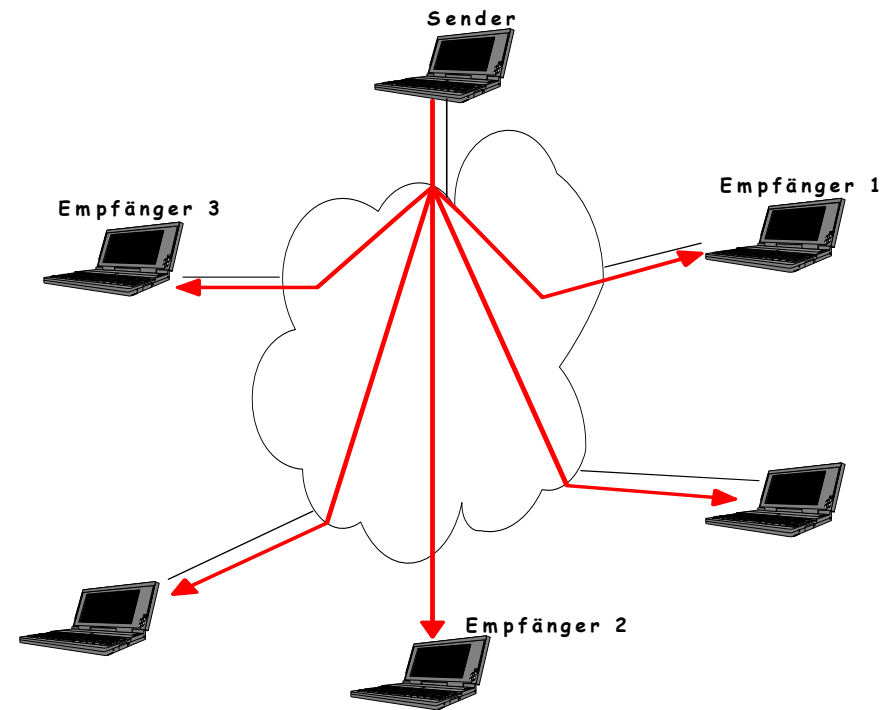
**⇒ Scalable Communication Paths needed to
Distribute Data in Parallel**

Ineffective Group Communication

Unicast

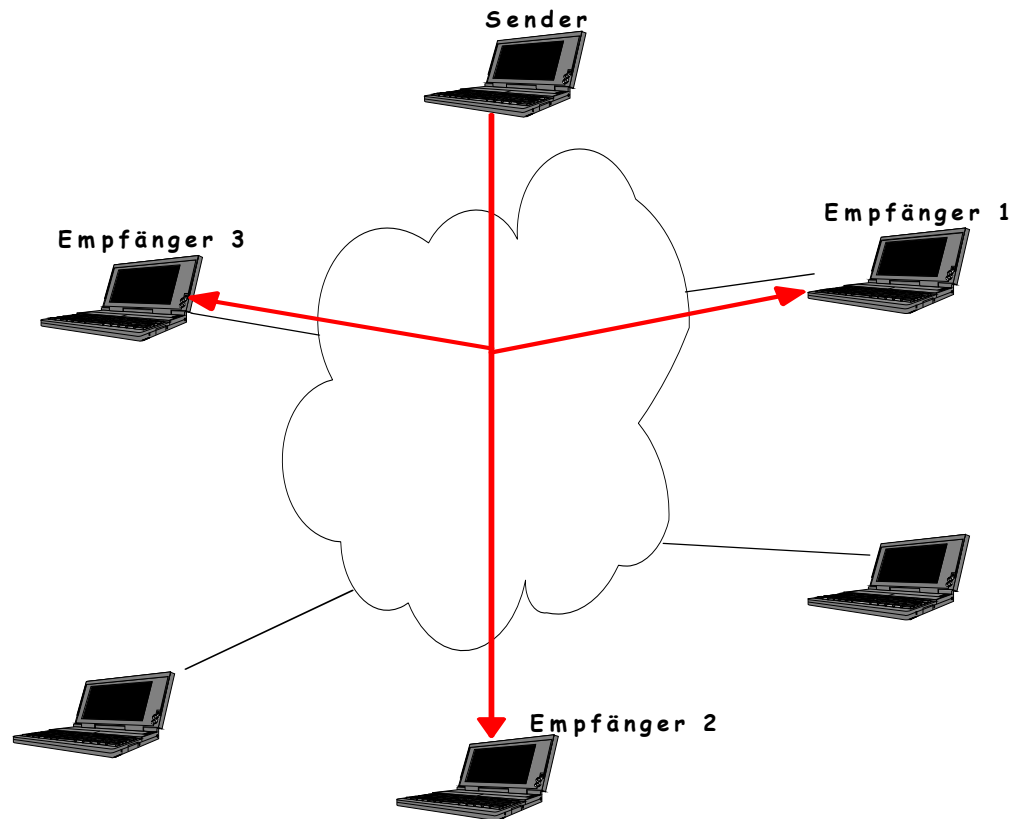


Broadcast



Effective Group Communication

Multicast



Group Communication Differs

Classical TCP/IP Communication Model:

- Client/Server Principle
- Individual Communication Channels
 - Initiated by client
 - Server answers individually
 - Server speaks on many point-to-point channels
- Exception for unspecific message distribution:
Broadcasts

Examples

IRC – Client-to-Client Communication via Server

NTP – many Clients ask one Server

Routing (RIPv1) – Broadcast of Routing Tables

Multisource Webpage – Client asks many Server

Internet Server Farm – one Client asks one of
many Servers

Group Communication Modes

Broadcast – one Sender to all Members of the Subnet

Concast – one Receiver of a Group of Senders

Multicast – one Sender Addresses a Group of Receivers

Multipeer – a Group of Senders to a Group of Receivers

Anycast – Communication Partner selected from a Group of potential Partners (Unicast)

Aspects of Group Communication

- **Openness:** Support of open and closed groups
- **Dynamic:** Change of group membership
- **Reliability:** Securing of data transport
- **Flow Control:** Adapting data streams to buffers
- **Group Management:** Mechanisms of addressing
and membership control

Openness & Dynamic

Relevant Mechanisms

- Identification/Announcement in a group/ with the sender
- Authorisation in closed Groups
- Management of send/receive allowances
- Registration & deregistration, definition of group composition
- Definition of group lifetime

Reliability

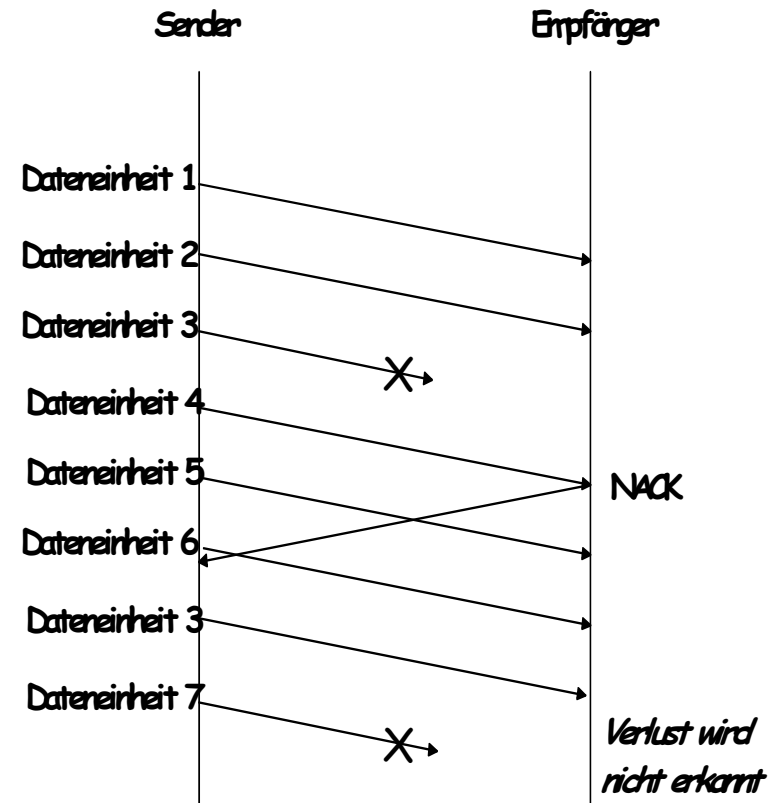
A securing layer requests for some acknowledgements

ACK:

- Group members need to register with sender
- results in ACK-'Implosion'

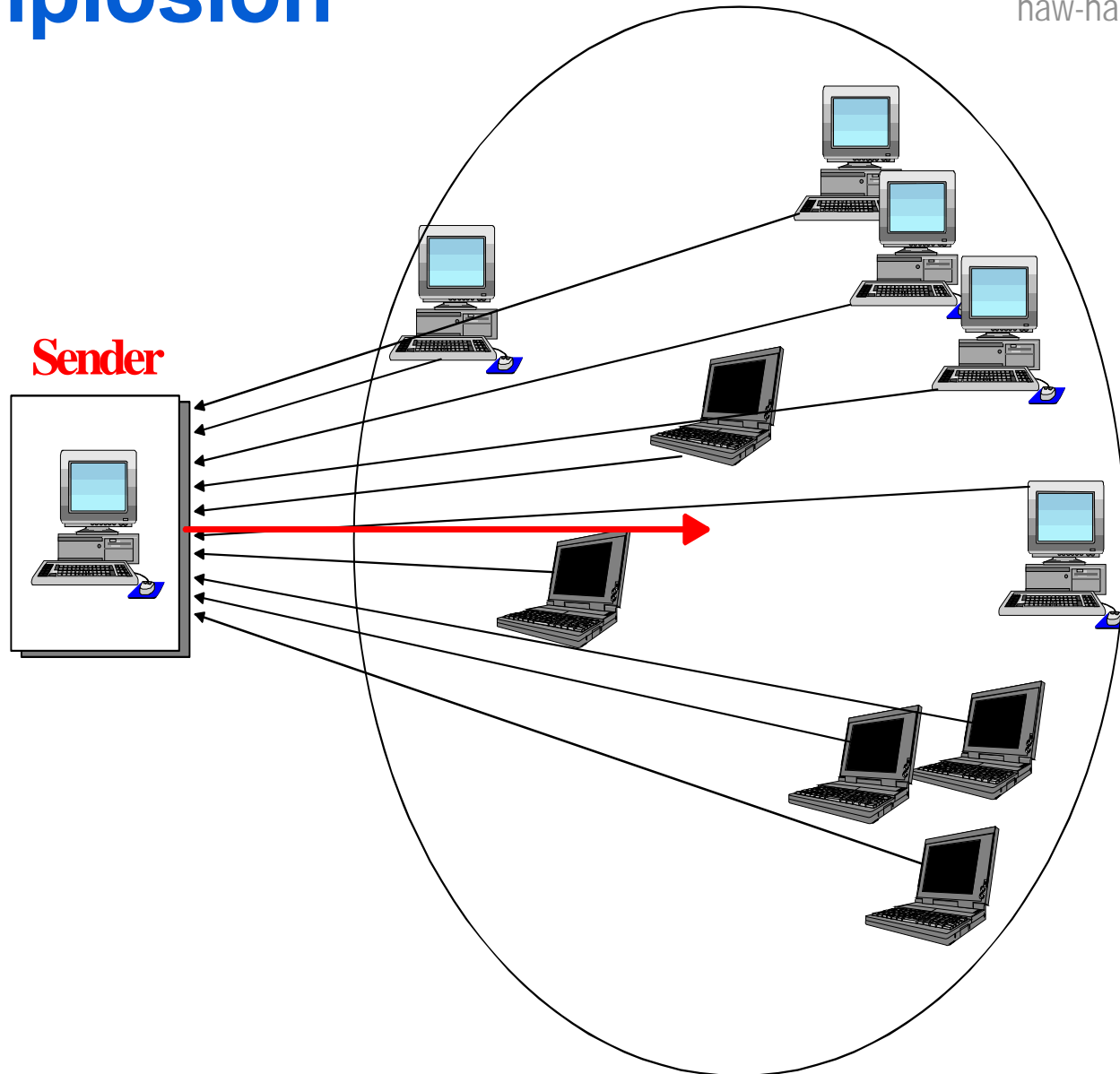
NACK:

- Retransmission for one may disturb the entire group
- Last loss may be unseen



Ack Implosion

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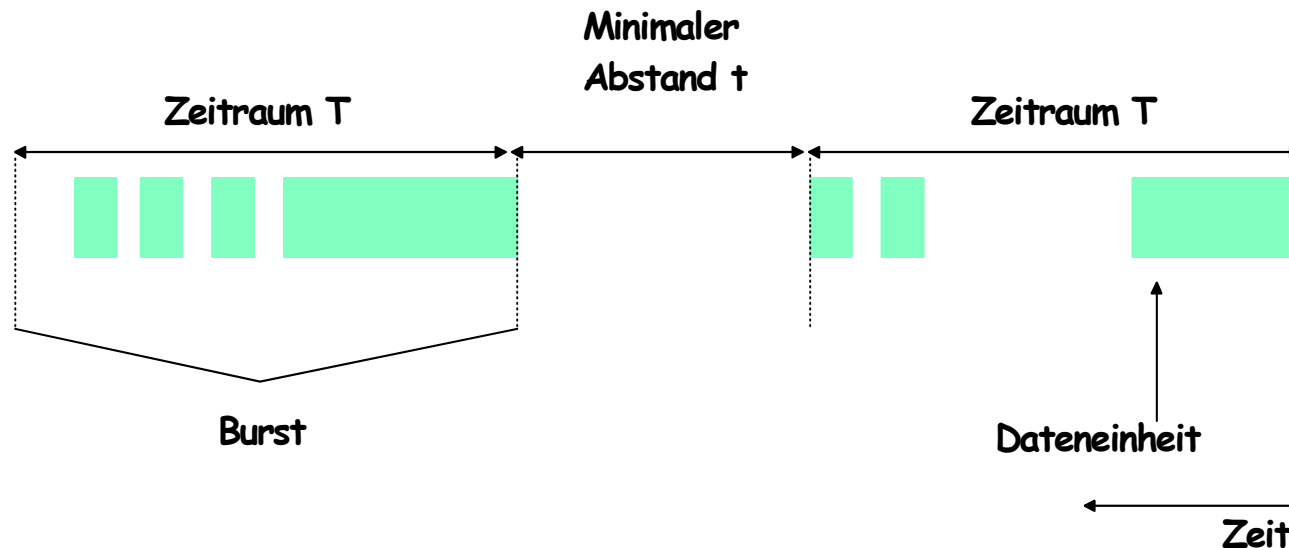
Flow Control

Window based:

- uses positive acknowledgements for sliding the window
⇒ unsuitable

Rate based:

- Adjust source intensity (Burst Rate)
- May be announced by receiver with membership registration



Group Management

Addressing

- Address scheme for a group
- Address allocation (centralised or decentralised)

Signaling

- Registration/Deregistration
- Member management (centralised or decentralised)

IP Multicasting

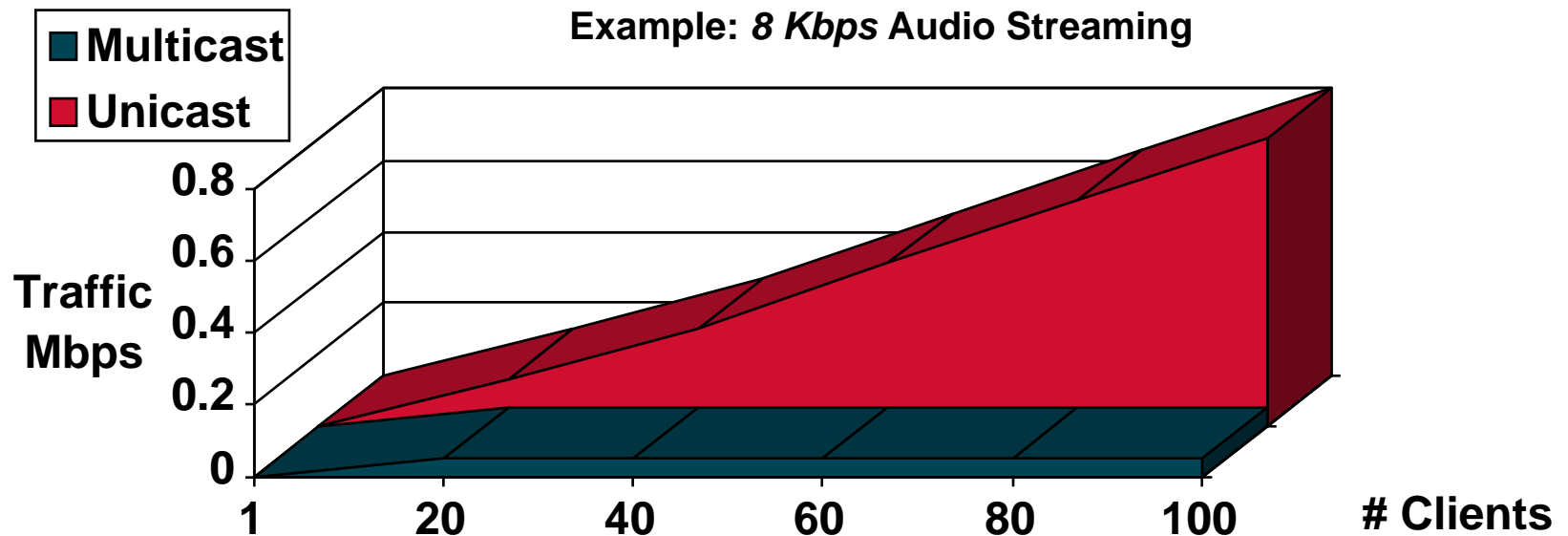
Method for Transferring IP Datagrams to Host-Groups

- Initially: RFC 1112 (S. Deering et.al., 1989)
- Addresses a host group by *one* group address
- Two kinds of multicast:
 - Any Source Multicast (ASM)
 - Source Specific Multicast (SSM)
- Client Protocol for registration (IGMP/MLD)
- Routing throughout the Internet (Multicast Routing)
- Address translation into Layer 2

Main Advantage of IP Multicasting

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- Omits redundant network traffic
- Reduces network and server load

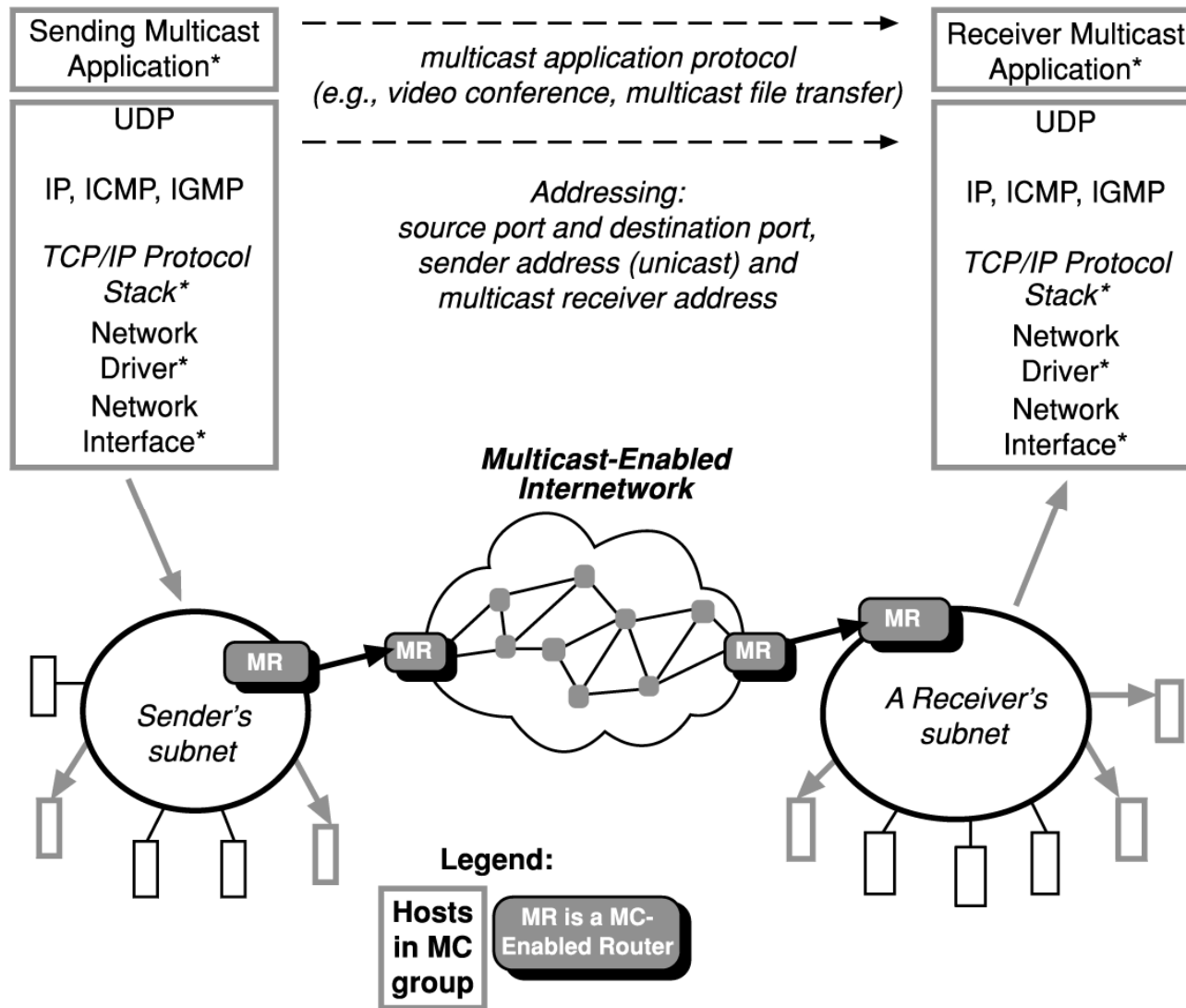


Aspects of IP Multicasting

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- Offers to sender a data delivery service to a distributed unknown group of receivers (multipoint access)
- UDP-based
- Best Effort Transport
- Securing and flow control left to application
- No closed groups
- No restriction on senders
- Applications may react source address sensitive

Multicast Network



Applications of IP Multicasting

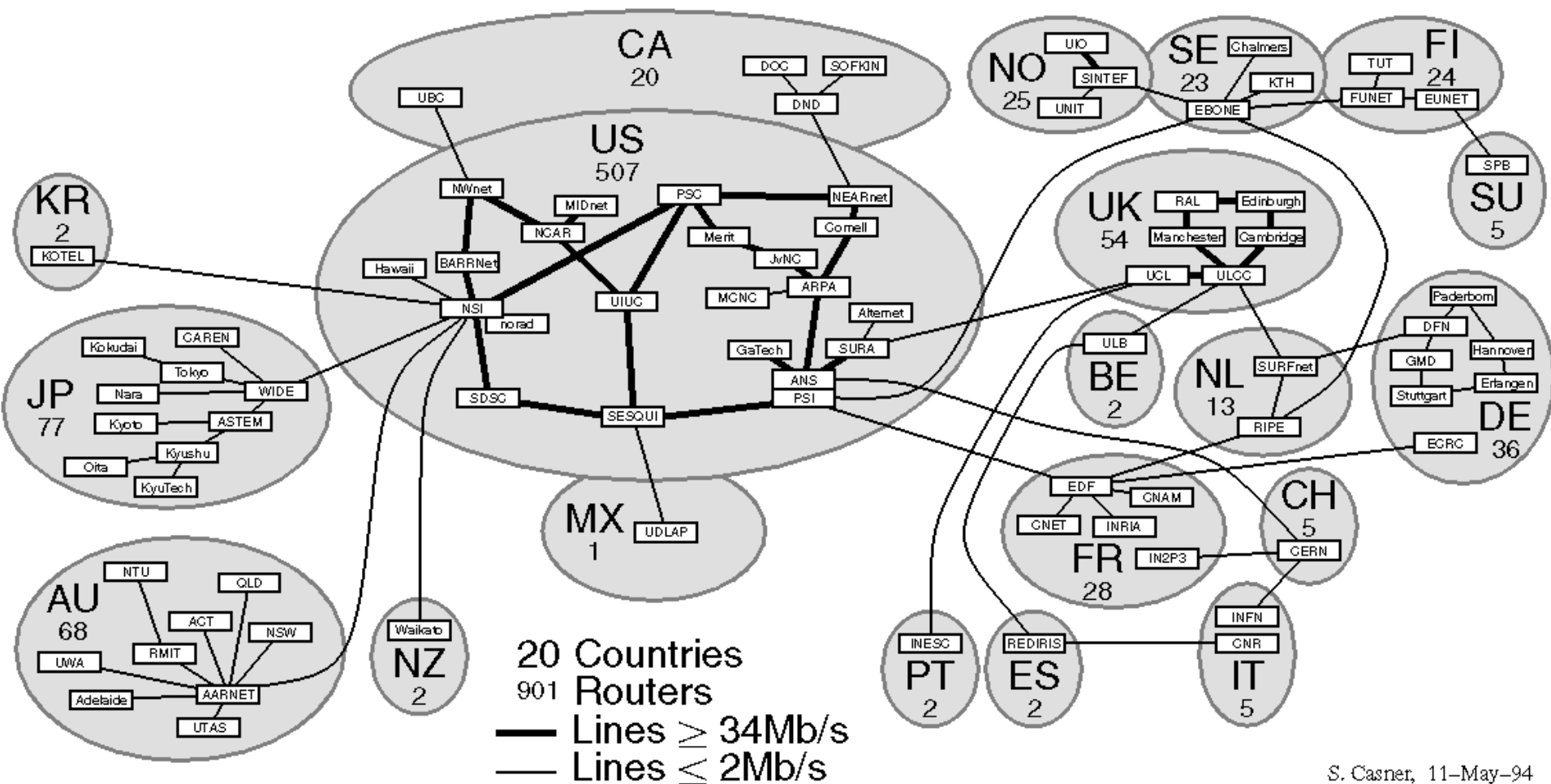
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- Multimedia
 - Streaming video and audio (broadcasting)
 - Teleteaching
 - Conferencing
- Financial information services (stock price ticker,...)
- Network information services
- Arbitrary data distribution services (Push Apps)

First Global Multicast Deployment: MBONE

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Major MBONE Routers and Links



Example:

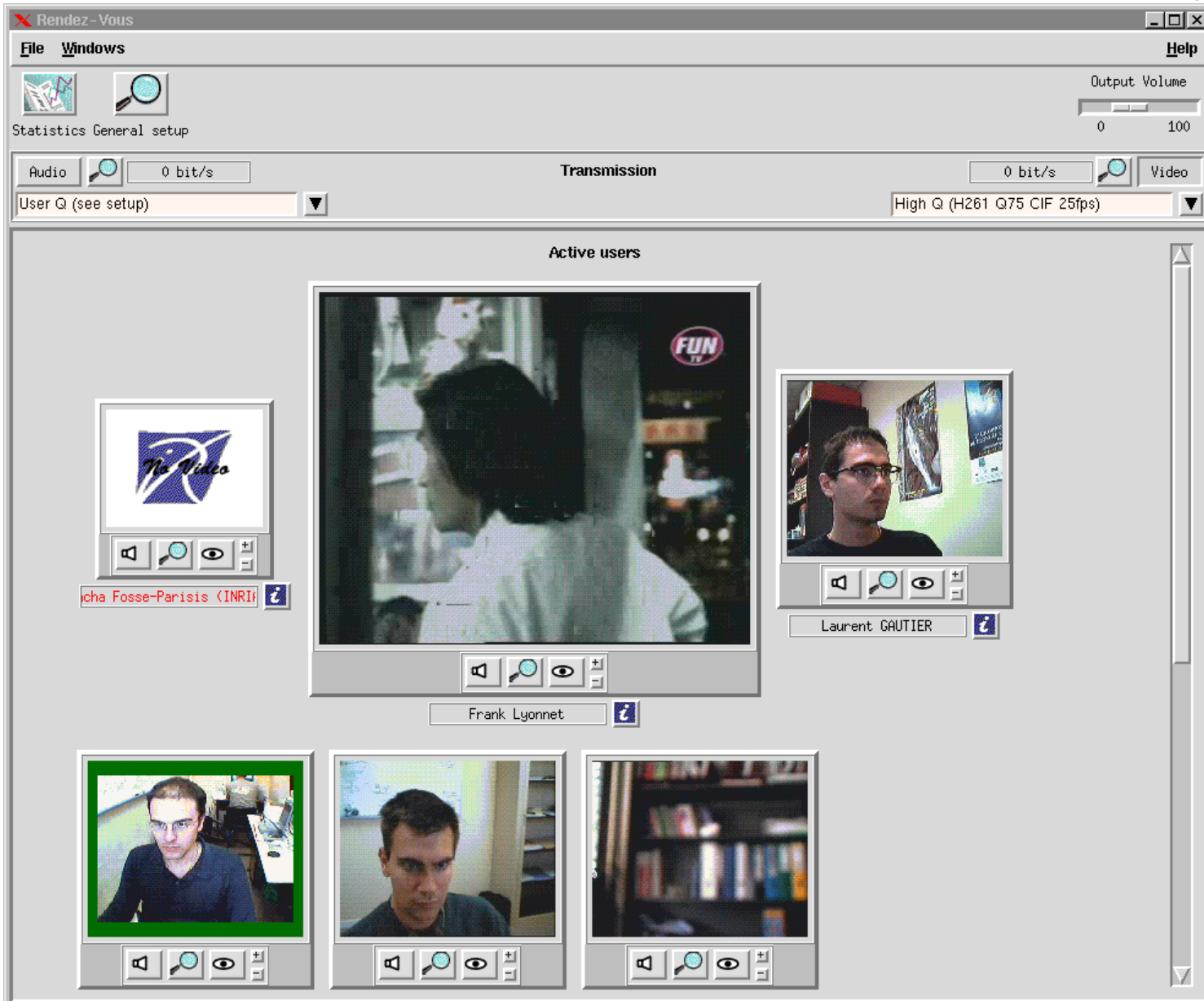
Mbone-Tools

SDR



Example: Mbone-Tools Rendez-Vous

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Multicast Addressing

- Denote delocalized addresses
- IPv4 Multicast Group addresses
 - 224.0.0.0–239.255.255.255
 - Class “D” Address Space
 - Special SSM block: 232.*.*.*
- IPv6: scoped multicast addresses
 - FF00::/8
 - Special SSM block: FF3x::/32
- Permanent Addresses assigned by IANA
 - RFC 1700: Assigned Addresses
 - *“<http://www.iana.org/assignments/multicast-addresses>”* lists reserved addresses
- Dynamic Addresses
 - independent of local IP-address space (IPv4)
 - Unicast based Multicast addresses (IPv6)

Internet Address Classes



=> max. 16.777.216 Hosts, IP-Adresse 1.x.y.z bis 127.x.y.z



=> max. 65.536 Hosts, IP-Adresse 128.x.y.z bis 191.x.y.z



=> max. 255 Hosts, IP-Adresse 192.x.y.z bis 223.x.y.z



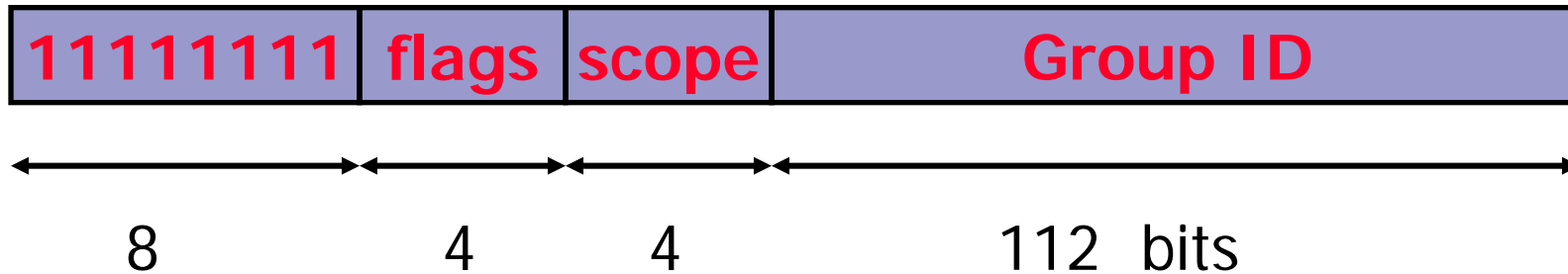
Private Multicast Addresses

- Officially not routed address range
 - 239.0.0.0–239.255.255.255
 - Private Address Space
 - Similar to RFC1918 Unicast Addresses
 - Unused for global Internet Traffic
 - Limits Multicast Traffic to own Institution
 - Same Addresses may be globally re-used
 - Example
 - Local range: 239.253.0.0/16
 - Organisation-wide range: 239.192.0.0/14

Reserved Multicast Addresses

- Permanent IP Multicast Group Addresses
 - 224.0.0.0–224.0.0.255
 - Examples:
 - 224.0.0.1 All Systems of Subnet
 - 224.0.0.2 All Routers of Subnet
 - 224.0.0.4 All DVMRP Router
 - 224.0.0.5 All OSPF Router
 - 224.0.0.9 All RIP(v2) Router
 - 224.0.0.13 All PIMv2 Router
 - 224.0.1.1 NTP
 - 224.0.1.9 Multicast Transport Protocol (MTP)
- TTL – Standards in MBONE
 - TTL = 1: This Subnet
 - TTL = 15: This Site
 - TTL = 63: This Region
 - TTL = 127: This Internet

IPv6 Multicast Addresses

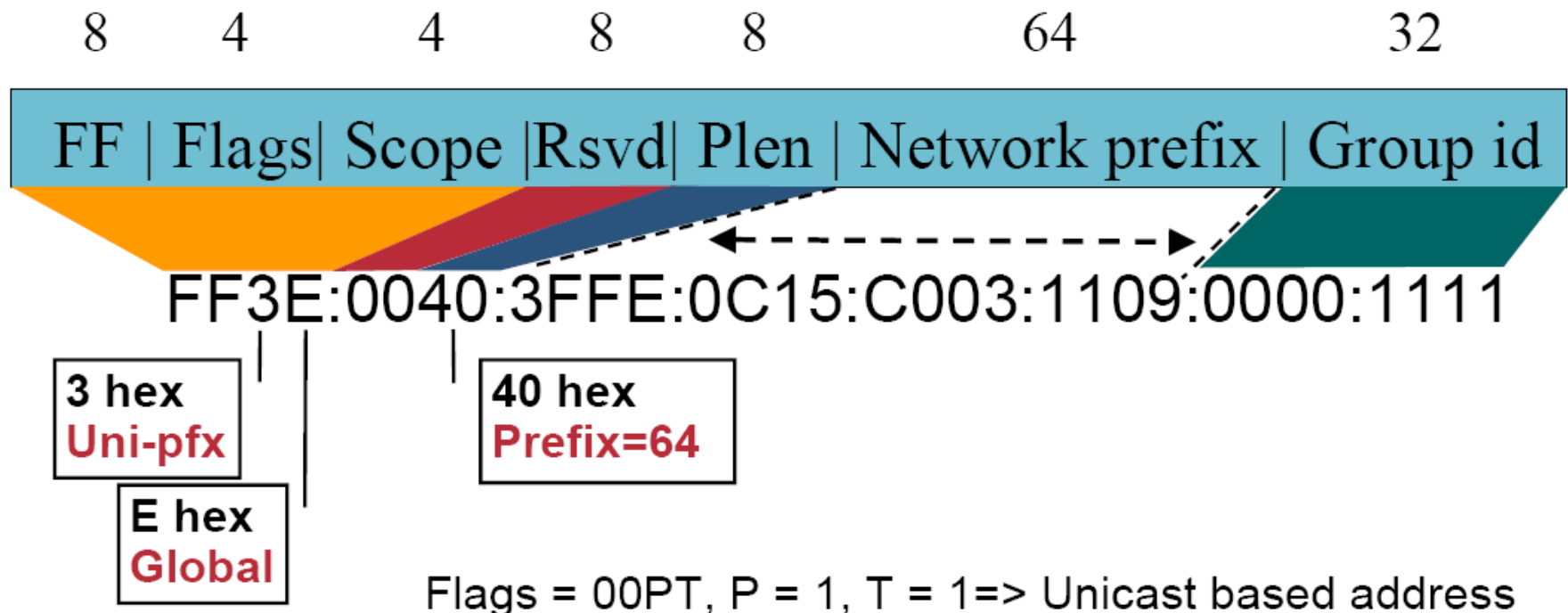


- **Flag field:** lower bit indicates permanent (=0) respectively transient (=1) group, rest is reserved (==0)
- **Scope field:**
 - 1 - node local
 - 2 - link-local
 - 5 - site-local
 - 8 - organisation local
 - B - community-local
 - E - global (other values reserved)

IPv6 Unicast Based Multicast Addresses (RFC 3306)

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- Solves the old IPv4 address assignment problem:
How can I get global IPv4 multicast addresses (GLOB, ..)
- In IPv6, if you own an IPv6 unicast address prefix you implicitly own an RFC3306 IPv6 multicast address prefix:



Dynamic Multicast Addressing

- Dynamic Assignment of Group addresses:
 - Until now: SDR Application
 - Sessions/Groups announced via well-known multicast groups
 - Address assignments and collisions are managed within initiation process
 - Brings up severe scaling issue
- Future Techniques and Planning:
 - Multicast Address Set-Claim (MASC)
 - Hierarchical, dynamical address allocation scheme
 - Difficult and far
 - MADCAP
 - Similar to DHCP
 - Needs own Protocol stack and application integration!

Internet Group Management

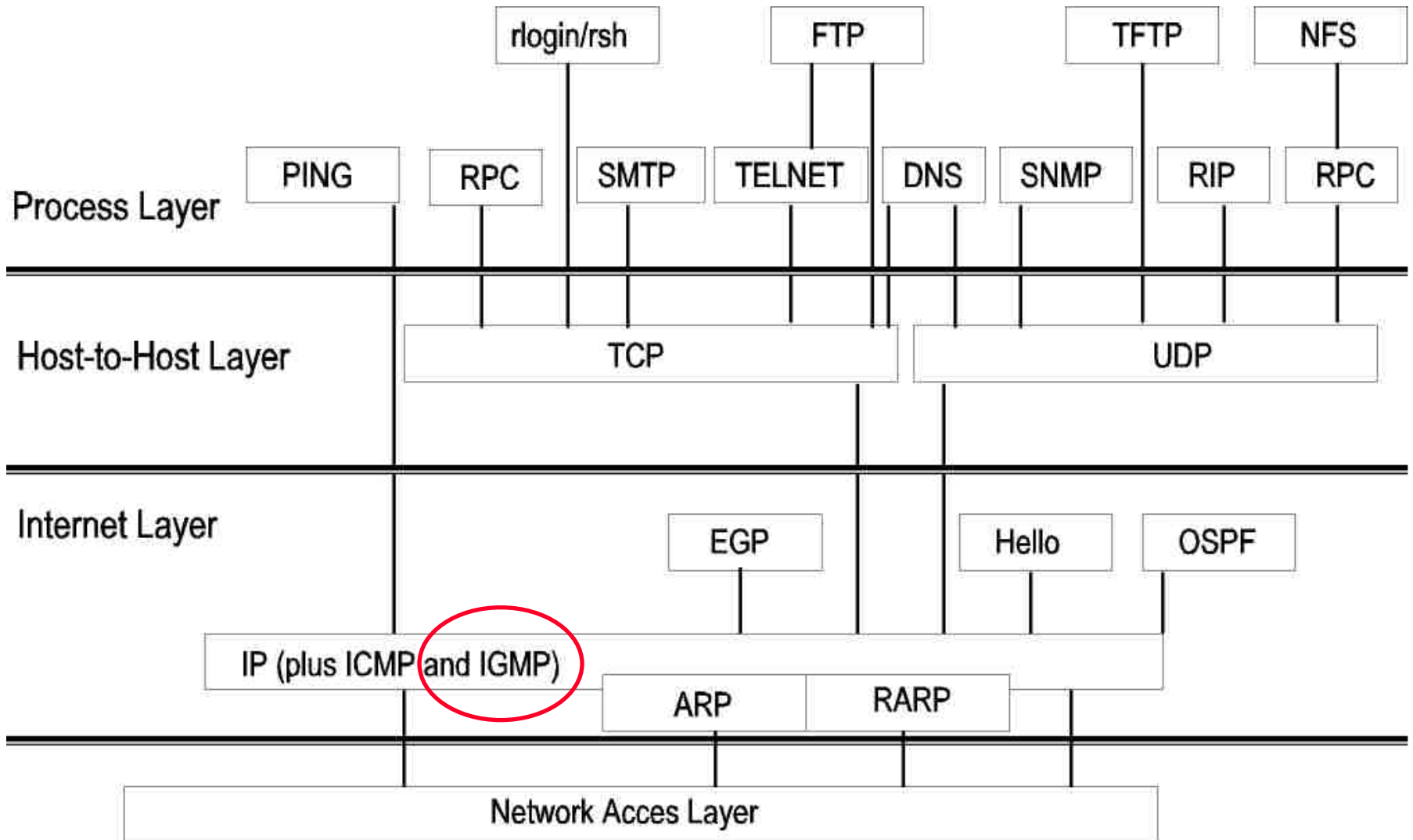
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Internet Group Management Protocol (IGMP)

- Client Protocol to initiate, preserve and terminate group membership
- Local Router collect and monitor information
- IPv4: Internet Group Management Protocol (IGMP)
 - IGMP v1 RFC 1112
 - IGMP v2 RFC 2236 – implemented almost everywhere
 - IGMP v3 RFC 3376
- IPv6: Multicast Listener Discovery Protocol (MLD)
 - MLDv1 (RFC 2710) – analogue to IGMPv2
 - MLDv2 (RFC 3810) – starting from IGMPv3
- SSM Specialities: RFC 4604

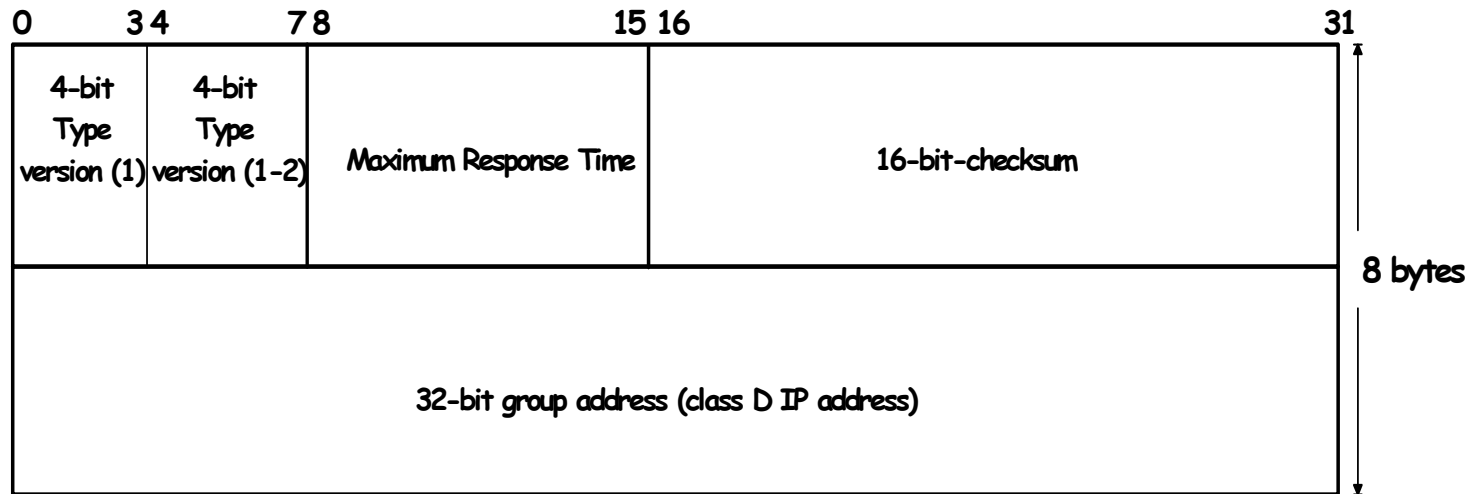
IGMP

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IGMP Protocol Architecture

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IGMP works like ICMP with Queries:

- General Membership
- Group specific Membership
- Version 2 Membership Report
- Leave Query
- Version 1 Membership Report

IGMP Kommunikation

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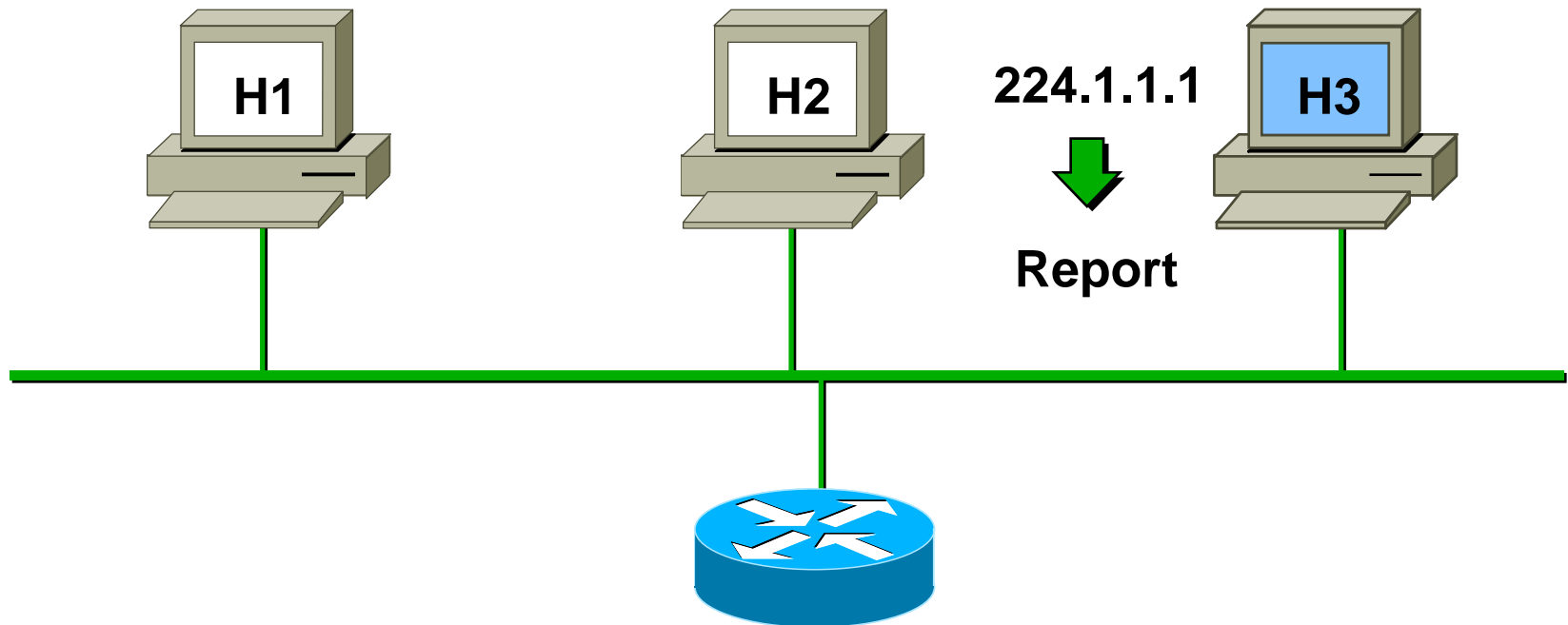
IGMP report, TTL = 1,
IGMP group addr = group address
dest IP addr = group address
src IP addr = host's IP addr

IGMP report, TTL = 1,
IGMP group addr = 0
dest IP addr = 224.0.0.1
src IP addr = router's IP addr



IGMP Host-Router Signalling

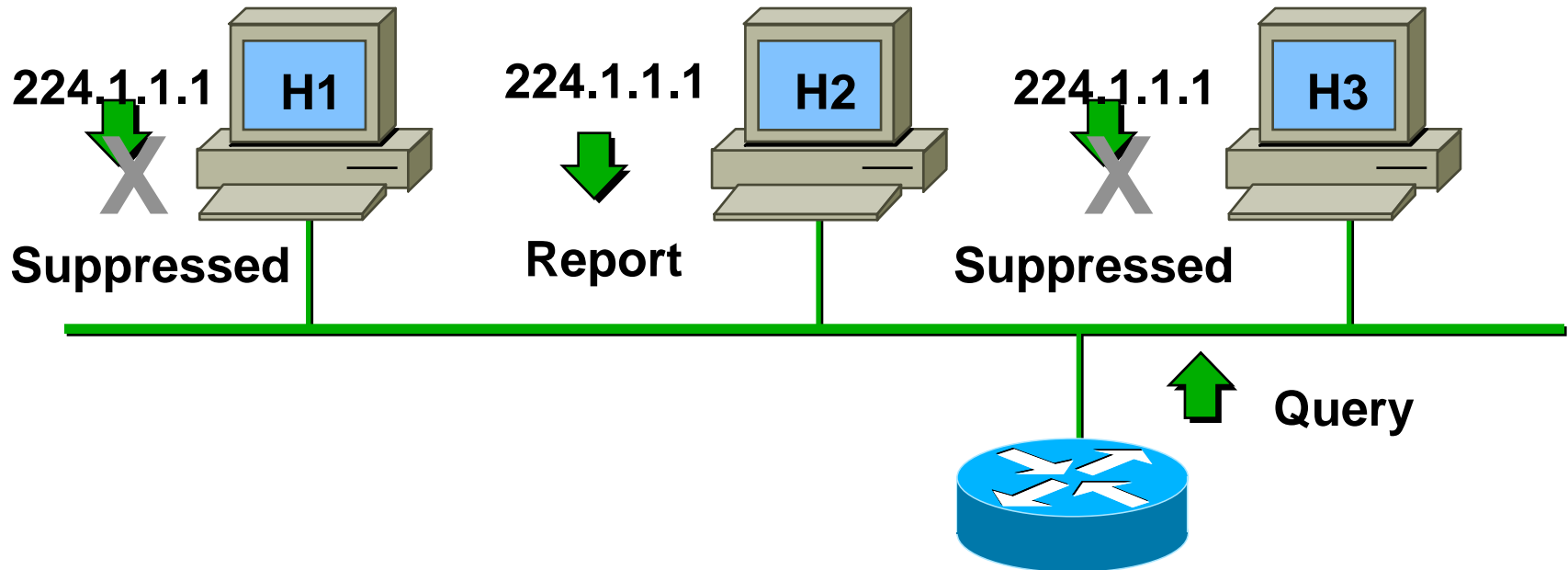
Group membership report



Members joining a group do not have to wait for a query to join. They send in an unsolicited report indicating their interest

IGMP Host-Router Signalling

Group Membership Preservation

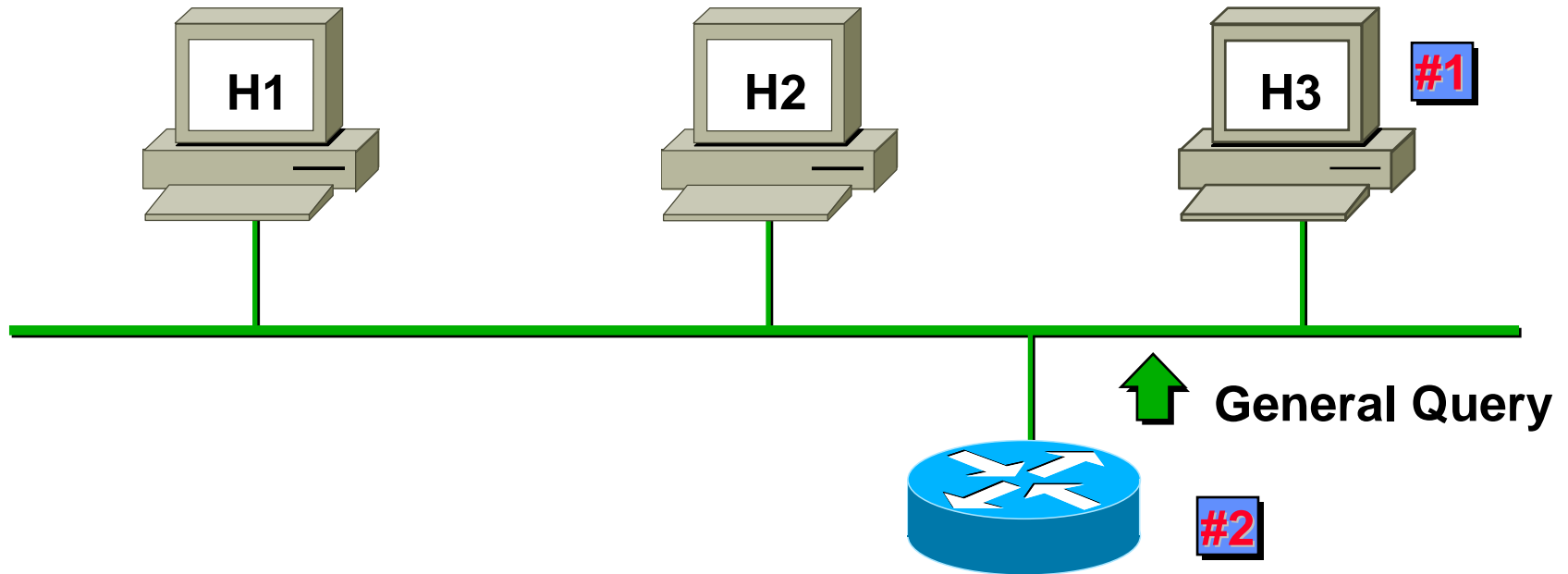


- Router sends periodic queries to 224.0.0.1
- One group member per subnet answers
- Others suppress answer

IGMP Host-Router Signalling

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Terminate Group Membership (IGMPv1)

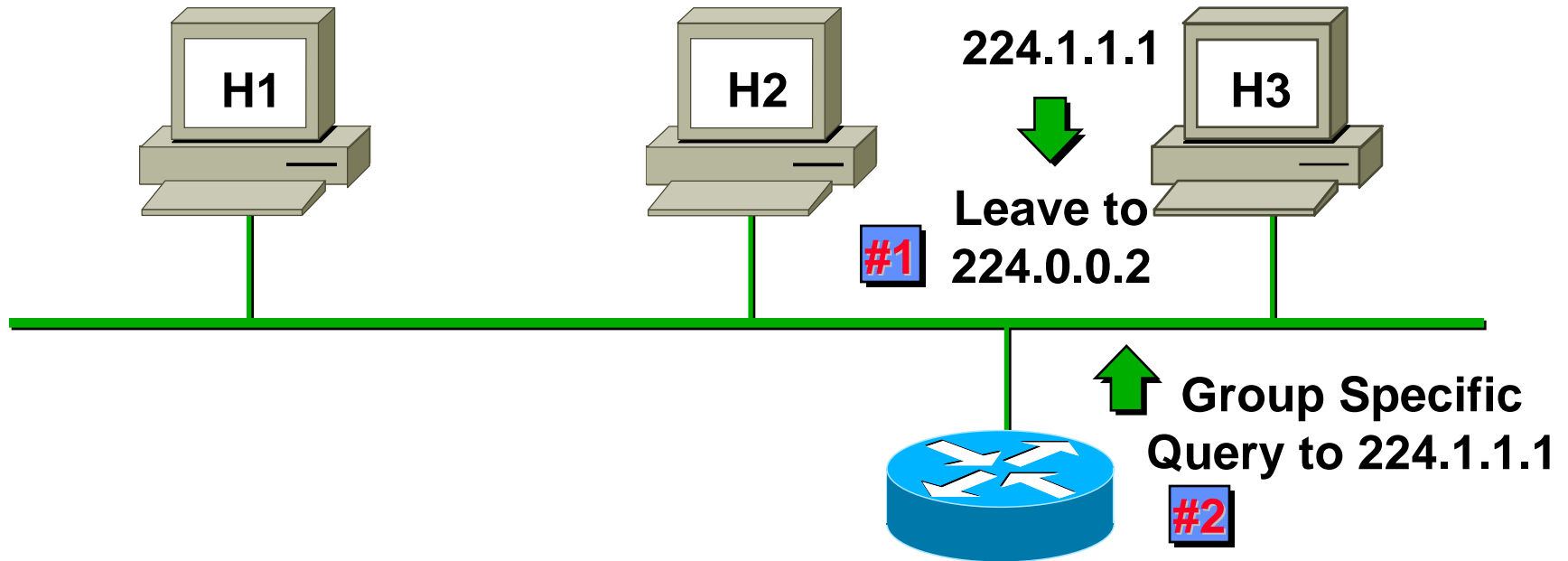


- Host 3 leaves group quietly
- Router queries remain unanswered
- Group terminate on timeout (up to 3 min)

IGMP Host-Router Signalling

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Terminate Group Membership (IGMPv2)



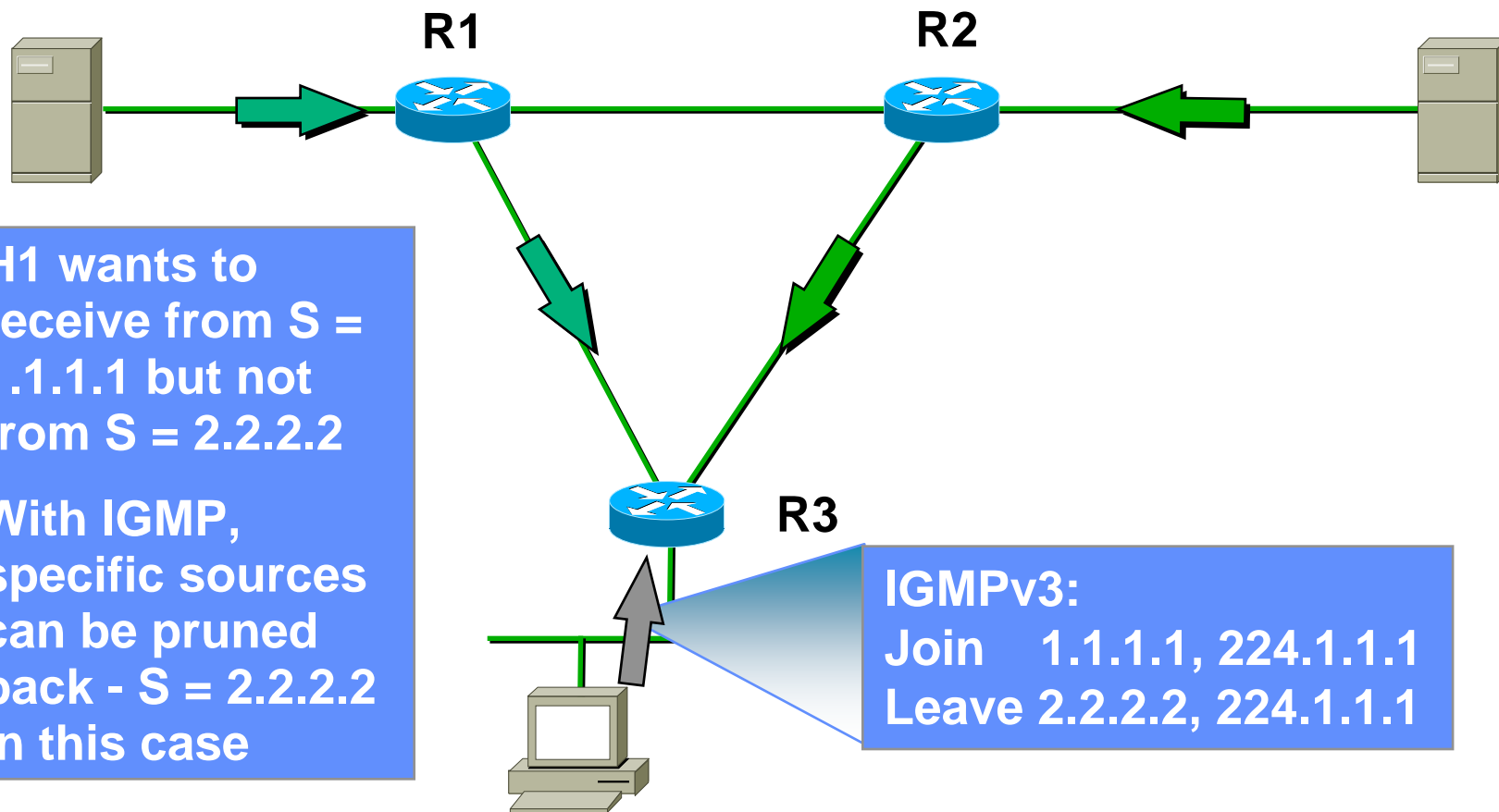
- Host sends Leave Message to 224.0.0.2
- Router sends group query to 224.1.1.1
- Timeout ~ 3 seconds for group 224.1.1.1

IGMP v3

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Source = 1.1.1.1
Group = 224.1.1.1

Source = 2.2.2.2
Group = 224.1.1.1



- H1 wants to receive from S = 1.1.1.1 but not from S = 2.2.2.2
- With IGMP, specific sources can be pruned back - S = 2.2.2.2 in this case

IGMPv3:
Join 1.1.1.1, 224.1.1.1
Leave 2.2.2.2, 224.1.1.1

H1 - Member of 224.1.1.1

Limits of IGMP

IGMP Concept has no Group Directory

- Hosts not answering on Membership Queries remain unseen
- Closed groups impossible
- Undiscovered listener part of the concept

IGMP is relatively slow

- Time to reaction in the order of seconds
- Unsuitable for flow control or congestion avoidance
- Initiation or change of a non-local group tardy

API

Berkeley Sockets set/getsockopt():

- **IP_ADD_MEMBERSHIP** *to join a multicast group on a specific interface*
- **IP_DROP_MEMBERSHIP** *to leave a multicast group (no protocol action initiated with IGMP v1, but there is with IGMP v2)*
- **IP_MULTICAST_IF** *to set or get default interface for use with multicast sends*
- **IP_MULTICAST_LOOP** *to disable loopback of outgoing multicast datagrams*
- **IP_MULTICAST_TTL** *to set the IP time-to-live of outgoing multicast datagrams.*

API - Java

Package: java.net

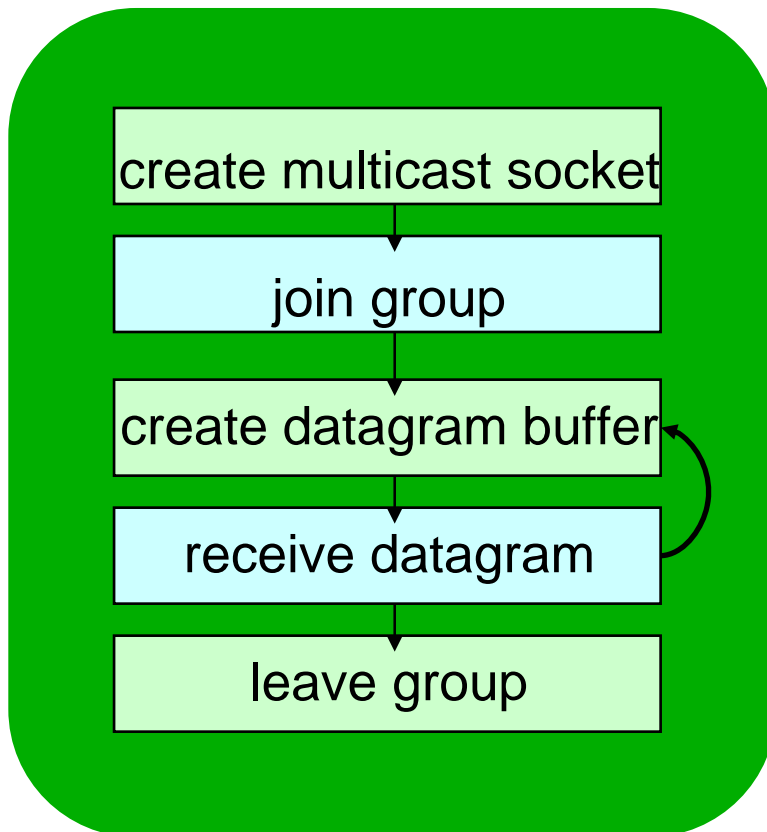
Class MulticastSocket

with Methods

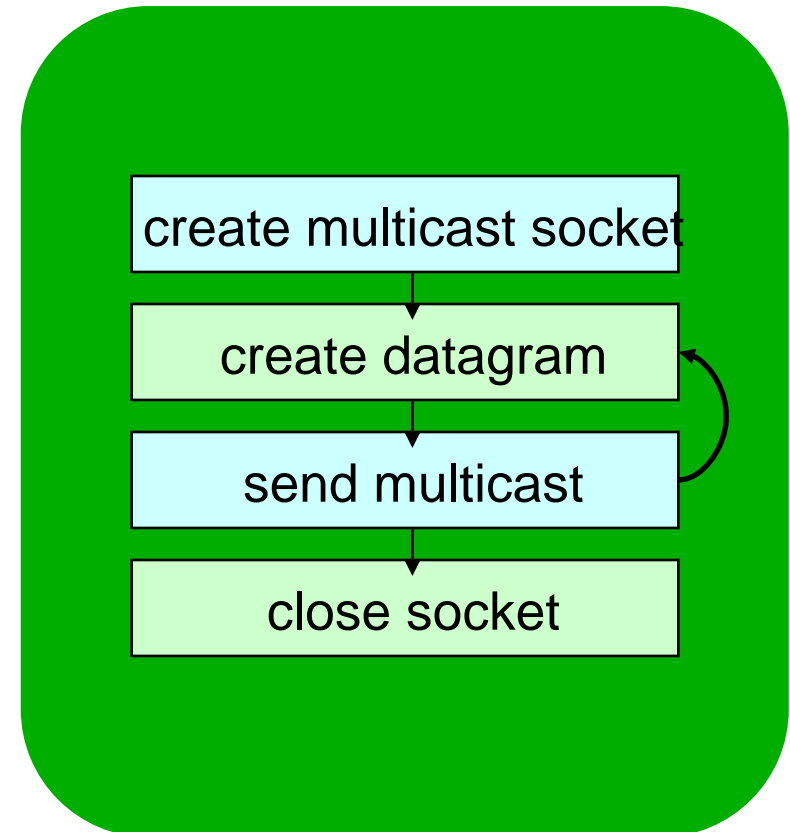
- *public void joinGroup(InetAddress mcastaddr)*
- *public void leaveGroup(InetAddress mcastaddr)*

IP Multicast in Java

Multicast Listener



Multicast Sender



IP Multicast in Java

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```
import java.net.*;
import java.io.*;
public class MulticastPeer{
    public static void main(String args[]){
        // args give message contents & destination multicast group
        // (e.g. "228.5.6.7")
        try {
            InetAddress group = InetAddress.getByName(args[1]);
            s = new MulticastSocket(6789);
            s.joinGroup(group);
            byte [] m = args[0].getBytes();
            DatagramPacket messageOut = new DatagramPacket(m, m.length, group, 6789);
            s.send(messageOut);
            // get messages from others in group
            byte[] buffer = new byte[1000];
            for(int i=0; i< 3; i++) {
                DatagramPacket messageIn = new DatagramPacket(buffer, buffer.length);
                s.receive(messageIn);
                System.out.println("Received:" + new String(messageIn.getData()));
            }
            s.leaveGroup(group);
        }catch (SocketException e){System.out.println("Socket: " +
e.getMessage());
        }catch (IOException e){System.out.println("IO: " + e.getMessage());}
    } } }
```

Reading

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- R. Wittmann, M. Zitterbart: Multicast Communication, Morgan Kaufmann, 2001